

NECESIDAD

en S_2 , S_3 , S_4 y S_5

Oscar A. Monroy

Si $\vdash_{s_2} \alpha$, entonces $\vdash_{s_3} \alpha$.

Si $\vdash_{s_4} \alpha$, entonces $\vdash_{s_5} \alpha$.

Paradojas de la implicación material

$$\not\vdash_{S3} Q \supset (P \supset Q)$$

$$\not\vdash_{S3} \neg P \supset (P \supset Q)$$

$$\not\vdash_{S3} ((P \wedge Q) \supset R) \supset ((P \supset R) \vee (Q \supset R))$$

$$\not\vdash_{S3} \neg(P \supset B) \supset P$$

Tampoco valen en S5.

S2

$\vdash_{S2} \Box a \rightarrow a$

$\vdash_{S2} a \rightarrow \Diamond a$

S3

$\vdash_{S3} \Box \alpha \rightarrow \alpha$

$\vdash_{S3} \alpha \rightarrow \Diamond \alpha$

$\vdash_{S3} (\alpha \rightarrow \beta) \rightarrow (\neg \Diamond \beta \rightarrow \neg \Diamond \alpha)$

S3

$\vdash_{S3} \alpha \rightarrow \alpha$, pero

$\not\vdash_{S3} \Box(\alpha \rightarrow \alpha)$

¿No es necesario que algo se implique a sí mismo?

S3

En S3 podemos mostrar que nuestras reglas lógicas son teoremas, pero no podemos mostrar que nuestras reglas lógicas son necesarias. ¿Qué tipo de necesidad es ésta? ¿Es la que queremos modelar?

S4

$$\vdash_{S4} \Box\alpha = \Box\Box\alpha$$

$$\vdash_{S4} \Diamond\alpha = \Diamond\Diamond\alpha$$

S4

$$\vdash_K \Box(a \rightarrow a)$$

S5

$$\vdash_{S5} \Box\alpha = \Diamond\Box\alpha$$

$$\vdash_{S5} \Diamond\alpha = \Box\Diamond\alpha$$

$$\vdash_{S5} \Box\alpha = \Box\Box\alpha$$

$$\vdash_{S5} \Diamond\alpha = \Diamond\Diamond\alpha$$

S5

$$\vdash_{S5} \Box \alpha = \Diamond \Box \alpha$$

$$\vdash_{S5} \Diamond \alpha = \Box \Diamond \alpha$$

$$\vdash_{S5} \Box \alpha = \Box \Box \alpha$$

$$\vdash_{S5} \Diamond \alpha = \Diamond \Diamond \alpha$$

S5

$$\square \diamond \square \square \diamond \square \alpha = ? ?$$

S5

$$\square \diamond \square \square \diamond \square \alpha = ?$$

$$\vdash_{S5} \square \alpha = \diamond \square \alpha$$

$$\vdash_{S5} \diamond \alpha = \square \diamond \alpha$$

$$\vdash_{S5} \square \alpha = \square \square \alpha$$

$$\vdash_{S5} \diamond \alpha = \diamond \diamond \alpha$$

S5

$$\square \diamond \square \square \diamond \square \alpha = ?$$

$$\vdash_{S5} \square \alpha = \diamond \square \alpha$$

$$\bullet \vdash_{S5} \diamond \alpha = \square \diamond \alpha$$

$$\vdash_{S5} \square \alpha = \square \square \alpha$$

$$\vdash_{S5} \diamond \alpha = \diamond \diamond \alpha$$

S5

$$\square \diamond \square \square \diamond \square \alpha = ?$$

● $\vdash_{S5} \square \alpha = \diamond \square \alpha$

$$\vdash_{S5} \diamond \alpha = \square \diamond \alpha$$

$$\vdash_{S5} \square \alpha = \square \square \alpha$$

$$\vdash_{S5} \diamond \alpha = \diamond \diamond \alpha$$

S5

$$\square \diamond \square \square \diamond \square \alpha = ?$$

$$\vdash_{S5} \square \alpha = \diamond \square \alpha$$

$$\vdash_{S5} \diamond \alpha = \square \diamond \alpha$$

$$\bullet \vdash_{S5} \square \alpha = \square \square \alpha$$

$$\vdash_{S5} \diamond \alpha = \diamond \diamond \alpha$$

S5

$$\square \diamond \square \square \diamond \square \alpha = ?$$

$$\vdash_{S5} \square \alpha = \diamond \square \alpha$$

$$\bullet \vdash_{S5} \diamond \alpha = \square \diamond \alpha$$

$$\vdash_{S5} \square \alpha = \square \square \alpha$$

$$\vdash_{S5} \diamond \alpha = \diamond \diamond \alpha$$

S5

$$\square \diamond \square \square \diamond \square \alpha = ?$$

● $\vdash_{S5} \square \alpha = \diamond \square \alpha$

$$\vdash_{S5} \diamond \alpha = \square \diamond \alpha$$

$$\vdash_{S5} \square \alpha = \square \square \alpha$$

$$\vdash_{S5} \diamond \alpha = \diamond \diamond \alpha$$

S5

$$\square \diamond \square \square \diamond \square \alpha = ?$$

$$\vdash_{S5} \square \alpha = \diamond \square \alpha$$

$$\vdash_{S5} \diamond \alpha = \square \diamond \alpha$$

$$\vdash_{S5} \square \alpha = \square \square \alpha$$

$$\vdash_{S5} \diamond \alpha = \diamond \diamond \alpha$$

S5

$$\square \diamond \square \square \diamond \square \alpha = \square \alpha$$

S5

$$\vdash_{S5} (\diamond a \supset \beta) = (a \supset \Box \beta)$$

$$\not\vdash_{S4} (\diamond a \supset \beta) = (a \supset \Box \beta)$$

¿Implicación en términos de necesidad?

$$\vdash_L \Box\beta \rightarrow (\alpha \rightarrow \beta)$$

$$\vdash_L \neg\Diamond\alpha \rightarrow (\alpha \rightarrow \beta)$$

Si $\vdash_L\alpha$, entonces $\vdash_{S2}\alpha$ y $\vdash_{S4}\alpha$.

¿Implicación en términos de necesidad?

$$\vdash_L \Box\beta \rightarrow (\alpha \rightarrow \beta)$$

$$\vdash_L \neg\Diamond\alpha \rightarrow (\alpha \rightarrow \beta)$$

$$\vdash_L \Box(\beta \vee \neg\beta)$$

$$\vdash_L \neg\Diamond(\alpha \wedge \neg\alpha)$$

Si $\vdash_L\alpha$, entonces $\vdash_{S2}\alpha$ y $\vdash_{S4}\alpha$.

¿Implicación en términos de necesidad?

$$\vdash_L \Box\beta \rightarrow (\alpha \rightarrow \beta)$$

$$\vdash_L \neg\Diamond\alpha \rightarrow (\alpha \rightarrow \beta)$$

$$\vdash_L \Box(\beta \vee \neg\beta)$$

$$\vdash_L \neg\Diamond(\alpha \wedge \neg\alpha)$$

$$\vdash_L \alpha \rightarrow (\beta \vee \neg\beta)$$

$$\vdash_L (\alpha \wedge \neg\alpha) \rightarrow \beta$$

Si $\vdash_L \alpha$, entonces $\vdash_{S2} \alpha$ y $\vdash_{S4} \alpha$.